

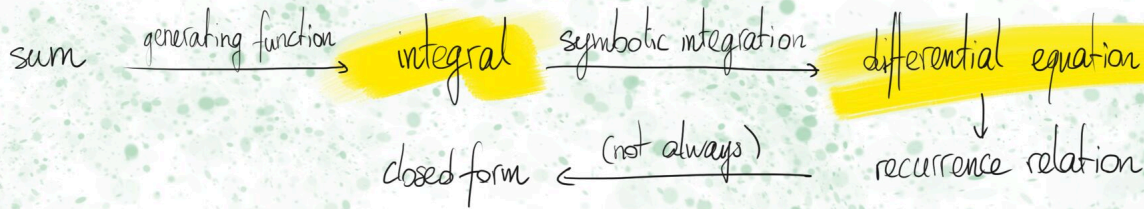
Computing with
integrals
in nonlinear algebra

by PIERRE LAIREZ

Discrete sums

$$\sum_{k=0}^{2n} (-1)^k \binom{2n}{k}^3 = (-1)^n \frac{(3n)!}{n!^3} \quad [\text{Dixon}]$$

$$\sum_{i,j \in \mathbb{Z}} \binom{2n}{n+i} \binom{2n}{n+j} |i^3 - j^3| = \frac{2n^2(5n-2)}{4n-1} \binom{4n}{2n} \quad [\text{Brent, Ohtsuka, Osborn, Prodinger}]$$



MOTIVATING EXAMPLE

Picard group

Find $a_0, \dots, a_4 \in \mathbb{C}^4$ st.

$$\forall t, f\left(\sum_{i=0}^4 a_i t^i\right) = 0$$

+ smoothness cond

$$X = \{x^4 + y^3z + xyzw + z^3w + yw^3 = 0\} \subset \mathbb{P}^3(\mathbb{C})$$

This surface contains: 0 lines, 0 conics, 0 twisted cubics
but 133 056 smooth rational curves of degree 4.

Obtained by computing integrals at high precision (~ 3000 digits),
and finding integer relations between them.

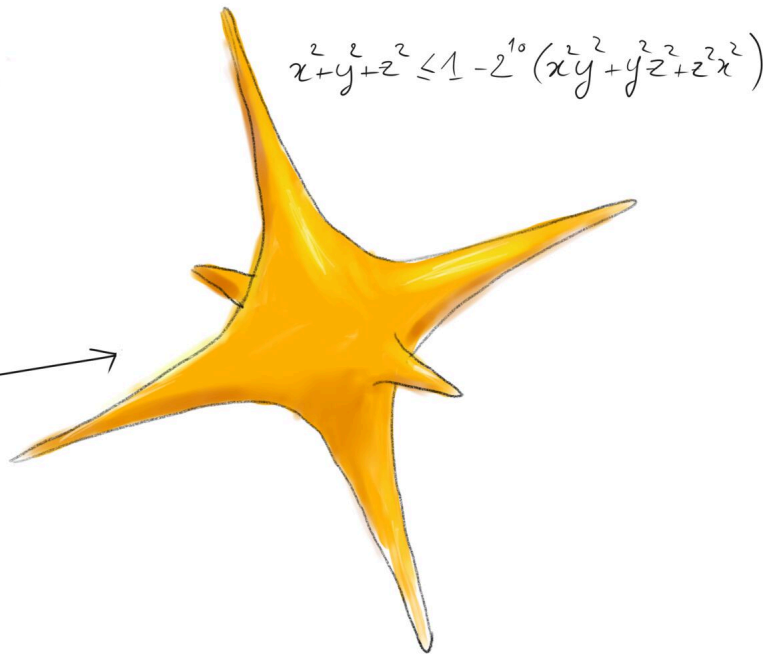
Challenge: compute all these curves!
(or just one)

Volume of semialgebraic sets

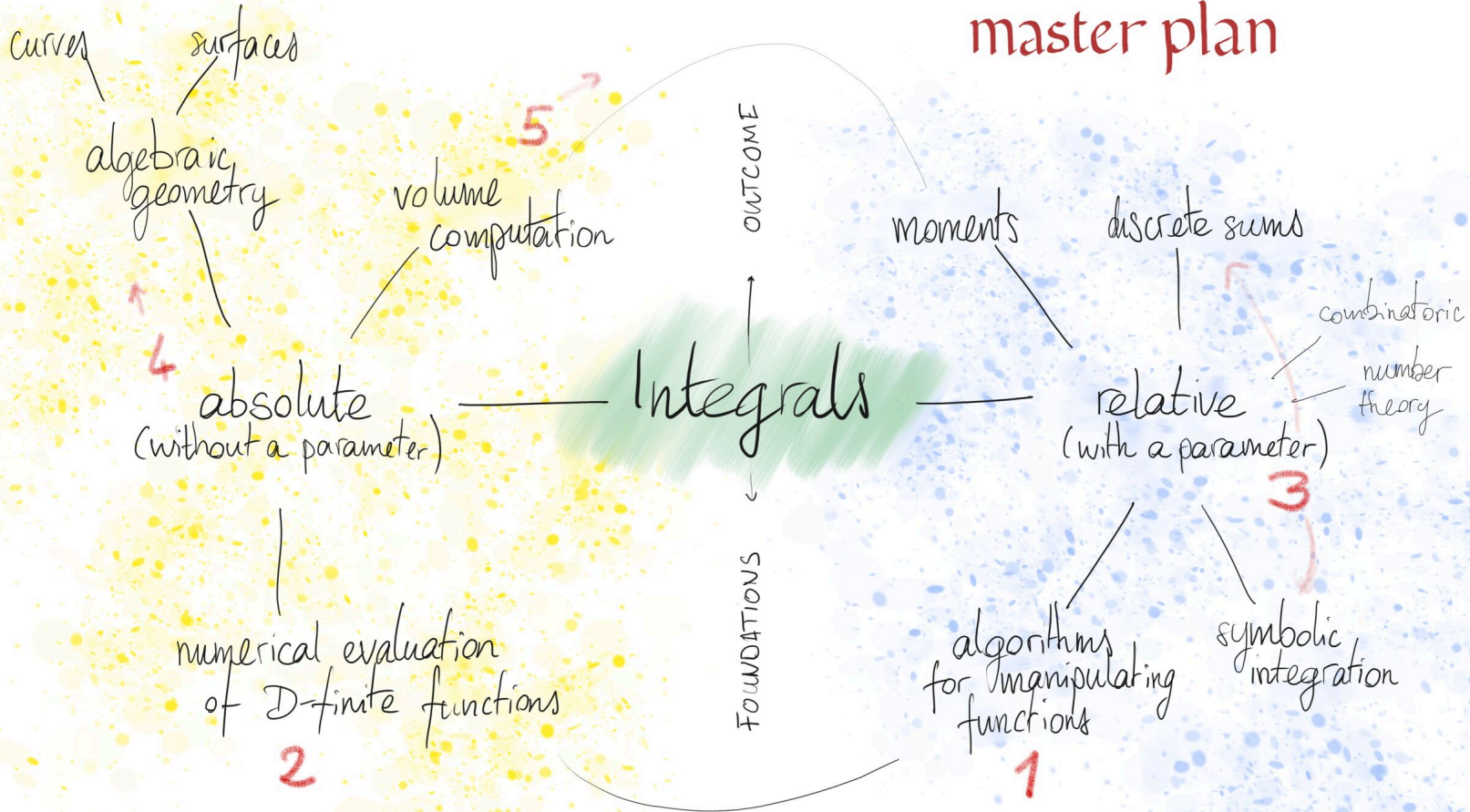
How to compute a thousand
digits of the volume of this?



Think of possible methods.
How will they scale?



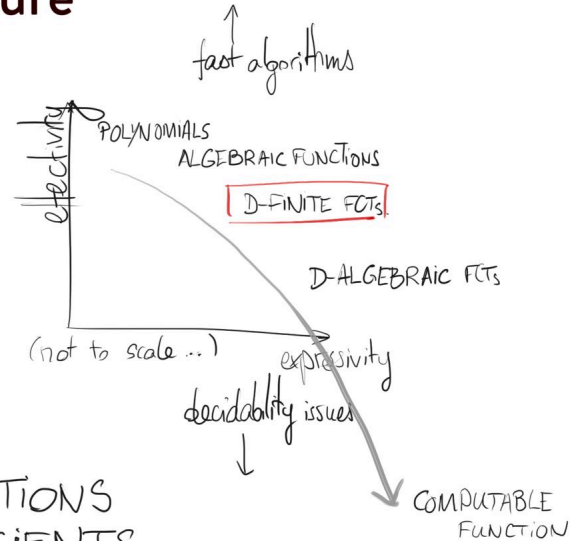
master plan



Differential equations as a data structure

To compute, we need representations
How to represent a function?
We want expressivity and effectivity
TRADE-OFF

LINEAR DIFFERENTIAL EQUATIONS
WITH POLYNOMIAL COEFFICIENTS
(+ initial conditions)

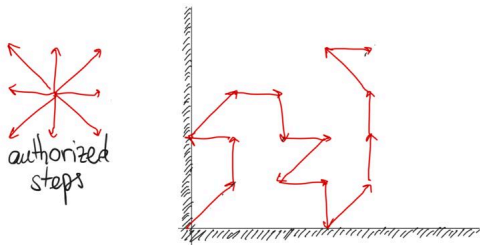


Elementary examples

$f(t)$	differential equation
1	$y' = 0$
t^n	$ty' = ny$
e^t	$y' = y$
$Ai(t)$	$y'' = ty$
$J_\alpha(t)$	$t^2 y'' + ty' + (t^2 - \alpha^2)y = 0$

Also many non-examples
 $\tan(t)$, $\Gamma(t)$, ...

Explicit vs. implicit



$a_n \stackrel{\text{def}}{=} \# \text{ of paths of length } n$
 $f(t) \stackrel{\text{def}}{=} \sum_{n \geq 0} a_n t^n$
(OEIS A151331)

Theorem (Bostan, Chyzak, van Hoeij, Kauers, Peh)

$$i. f(t) = \frac{1}{t} \int_0^t \frac{1}{(1+4u)^3} {}_2F_1\left(\frac{3}{2}, \frac{3}{2} \middle| \frac{16u(1+u)}{(1+4u)^2}\right) du$$

$$ii. (64t^6 + 40t^5 - 30t^4 - 5t^3 + t^2) f''' + (\dots) f'' + \dots = 0$$

EXPLICIT: reveals the nature of the function
does not always exist

IMPLICIT: enable computations

The algebra of differential operators

$\mathcal{D} \equiv \mathbb{C}[t]\langle \partial \rangle$ is the algebra generated by t and ∂
($\mathcal{D}_{\text{rat}} \equiv \mathbb{C}(t)\langle \partial \rangle$) subject to the relation $\partial t = t\partial + 1$
(Leibniz rule)
It acts on "functions"

For $L = \sum_{i=0}^r p_i(t) \partial^i \in \mathcal{D}$ and f a "function", $L \cdot f : t \mapsto \sum_{i=0}^r p_i(t) f^{(i)}(t)$
(normal form of an operator)

Examples: $\partial \cdot f = f'$
 $p(t) \cdot f = p(t)f$

$$(\partial - 1) \cdot \exp = 0$$
$$(t\partial - n) \cdot t^n = 0$$

diff. op. \equiv linear diff. eq.

Basic properties of differential operators

For $L \in \mathcal{D}_{\text{rat}}$, $\text{ord } L =$ maximum exponent of ∂ in L .

Lemma. (Euclidean division).

$\forall A, B \in \mathcal{D}_{\text{rat}}, A \neq 0 \Rightarrow \exists! Q, R \in \mathcal{D}_{\text{rat}} \quad A = QB + R \quad \text{and} \quad \text{ord } R < \text{ord } A$

that's actually a right division and there is a similar left division

Corollary: \mathcal{D}_{rat} is a principal ideal domain.

We define the right G.C.D. by

$$\mathcal{D}_{\text{rat}} \text{ rgcd}(A, B) = \mathcal{D}_{\text{rat}} A + \mathcal{D}_{\text{rat}} B$$

LCM by

$$\mathcal{D}_{\text{rat}} \text{ r lcm}(A, B) = \mathcal{D}_{\text{rat}} A \cap \mathcal{D}_{\text{rat}} B$$

the minimal eq. of f by

$$\mathcal{D}_{\text{rat}} \text{ ann}(f) = \{ A \in \mathcal{D}_{\text{rat}} \mid A \cdot f = 0 \}$$

LIFT ideals

Differential operators determine analytic functions

Theorem. Let $L \in \mathcal{D}$ and $a_0, \dots, a_{r-1} \in \mathbb{C}$ this is where the leading coefficient of L vanishes

For all simply connected open set $U \subseteq \mathbb{C} \setminus \text{sing } L$ and all $z \in U$ there is a unique holomorphic function f on U s.t.

$$L \cdot f = 0 \quad \text{and} \quad f^{(i)}(z) = a_i \quad (0 \leq i < \text{ord } L)$$

A function is represented by a differential operator and initial conditions
Representable functions are called **differentially finite** (at an ordinary point)
(or **D-finite**)

Structure of differentially finite power series

Let $L = \sum_{i=0}^r p_i(t) \partial^i \in \mathcal{D}$ and $f = \sum_{n=0}^{+\infty} a_n t^n \in \mathbb{C}[[t]]$ st $L \cdot f = 0$

Let $\theta = t\partial$ (Euler's derivation) and write $L = \sum_{i=0}^r \tilde{p}_i(t) \theta^i$
(NB: $\theta \cdot t^n = n t^n$)
 $= \sum_{j=0}^d t^j q_j(\theta)$

$$L \cdot f = \sum_{n=0}^{+\infty} \sum_{j=0}^d t^j q_j(\theta) \cdot a_n t^n = \sum_n \sum_j a_n q_j(n) t^{n+j} = \sum_{n=0}^{+\infty} t^n \left(\sum_{j=0}^{\min(n,d)} q_j(n-j) a_{n-j} \right)$$

indicial
polynomial
at 0

$$\underline{q_0(n) a_n = -q_1(n-1) a_{n-1} - \dots - q_d(n-d) a_{n-d}}$$

The sequence of coefficients is polynomially recursive (P-recursive)

Sequences and generating functions

$$(a_n)_{n \geq 0}$$

$$a_{n+1}$$

$$n a_n$$

P-recursive

asymptotic behavior

$$f = \sum_{n=0}^{\infty} a_n t^n$$

$$\frac{1}{t} f$$

$$t \partial \cdot f$$

D-finite

position and nature of the poles of f

Exercises

#1 Compute the coefficient of t^{1000}
 in $(1+t)^{10000} (1+t+t^3)^{50000}$

#2 A Motzkin word is a word on the alphabet $\{ '(', ')', 'x' \}$
 where each opening parenthesis is closed and each closing parenthesis
 closes an opening one. $M_n \stackrel{\text{def}}{=} \# \text{ of Motzkin word of length } n$

$M_4 = 9$ $(()), (xx), ()(), ()xx, (x)x, x()x, x(x), xx(), xxxx$

Show that $(n+4)M_{n+2} = (5+2n)M_{n+1} + (3+3n)M_n$

Guessing differential equations

$$\sum_{k=0}^n \binom{n}{k}^3 \quad (\text{Franel numbers})$$

1, 2, 10, 56, 346, 2252, 15184, 104960, 739162, 5280932, 38165260, 278415920, 2046924400, ???

$$f(t) = 1 + 2t + 10t^2 + 56t^3 + \dots + O(t^{13})$$

Search for polynomials $p_0(t), \dots, p_r(t)$ such that

$$p_0(t) f(t) + p_1(t) t f'(t) + \dots + p_r(t) t^r f^{(r)}(t) = O(t^{13})$$

\rightsquigarrow linear system of 13 equations in $\sum_i \deg p_i + 1$ unknowns.

We can try $r=2$ and $\deg p_i = 2$.

$$\rightsquigarrow t(t+1)(8t-1)\partial^2 + (24t^2+14t-1)\partial + (8t+2)$$

Extremely useful

Reliable

Proves nothing

 Maple

Diagonals of rational functions

in practice, this
will be \mathbb{Q}

$$\text{Let } F(z_1, \dots, z_n) = \frac{A}{P} \in \mathbb{C}(z_0, \dots, z_n) \quad P(0, \dots, 0) \neq 0$$

$$\text{Expand } F \text{ as } F = \sum_{k_0, \dots, k_n \geq 0} a_{k_0, \dots, k_n} z_0^{k_0} \dots z_n^{k_n} \in \mathbb{C}[[z_0, \dots, z_n]]$$

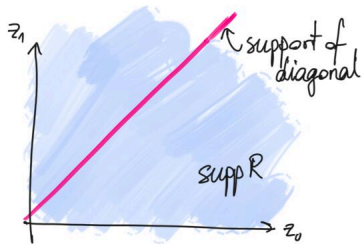
$$\text{Define } \text{diag} F \stackrel{\text{def}}{=} \sum_{k \geq 0} a_{k, \dots, k} t^k \in \mathbb{C}[[t]]$$

Theorem (Christol 1985) The power series $\text{diag} F$ is differentially finite

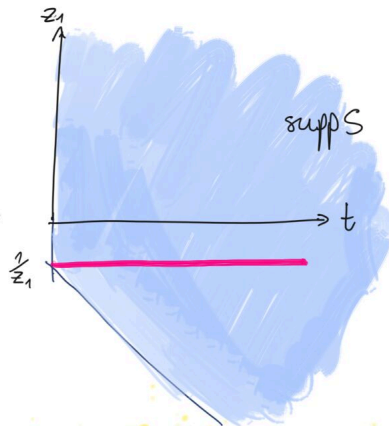
proof on board

A computational handle on diagonals

$$R(z_0, \dots, z_n)$$



$$S(t, z_1, \dots, z_n) = \frac{1}{z_1 \cdots z_n} R\left(\frac{t}{z_1 \cdots z_n}, z_1, \dots, z_n\right)$$



$$\mathbb{C}[z_1, \dots, z_n] \stackrel{(\ast)}{\left[\frac{1}{z_1 \cdots z_n} \right]} \llbracket t \rrbracket$$

notation for coefficient extraction

$$\text{res } S \stackrel{\text{def}}{=} \sum_{k \geq 0} \left(\left[\frac{t^k}{z_1 \cdots z_n} \right] S \right) t^k$$

Lemma. $\cdot \text{diag } R = \text{res } S$

$$\cdot \mathcal{L}(t, \partial_t) \cdot \text{diag } R = \text{res}(\mathcal{L} \cdot S)$$

$$\cdot \text{res} \left(\sum_{i=1}^n \frac{\partial S_i}{\partial z_i} \right) = 0$$

Mathematica

Corollary. If there are $A_1, \dots, A_n \in \mathbb{C}(t, z_1, \dots, z_n)$ and $\mathcal{L} \in \mathcal{D}$ st. $\mathcal{L} \cdot S = \frac{\partial A_1}{\partial z_1} + \dots + \frac{\partial A_n}{\partial z_n}$, then $\mathcal{L} \cdot \text{diag } R = 0$